An algebraic approach to Vaught's conjecture

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Tarskian Algebraic Logic – an area interdisciplinary between logic, and algebra (in fact the natural interface between universal algebra and logic) with an accompanying extremely rich geometry that has a varying dimension possibly transfinite-reflected in Tarski's cylindric algebras now better known as Concept Algebras when applied to (the algebraization of) sophisticated first order theories like spacetime geometries.



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The canonical examples of the so-called representable algebras, the cylindric set algebra, provide a natural vehicle for Model Theory, since cylindrifications reflect the semantics of existential quantifiers in logic, and are simply forming cylinders that is to say projections in Geometry. Using cylindric set algebras we approach Vaught's conjecture.

In 1961, Robert Vaught asked the following question: Given a complete theory in a countable language, is it the case that it either has countably many or 2^{\aleph_0} non-isomorphic countable models? By the number of non-isomorphic countable models is meant the number of their isomorphism-types; that is the number of equivalence classes of countable models w.r.t. the isomorphism relation between structures. We shall just say "the number of countable models" to mean the number of their isomorphism-types. In 1961, Robert Vaught asked the following question: Given a complete theory in a countable language, is it the case that it either has countably many or 2^{\aleph_0} non-isomorphic countable models? By the number of non-isomorphic countable models is meant the number of their isomorphism-types; that is the number of equivalence classes of countable models w.r.t. the isomorphism relation between structures. We shall just say "the number of countable models" to mean the number of their isomorphism-types.

The positive answer to the question is more commonly know as Vaught's Conjecture. (Vaught's conjecture has the reputation of being the most important open problem in model theory.) However, some logicians do not agree to this sweeping statement.

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Morley proved that the number of countable models is either less than or equal to the first uncountable cardinal ($\leq \aleph_1$) or else it has the power of the continuum. This is the best known (general) answer to Vaught's question.

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- 5. (Steel)theories of trees.

There are also attempts concerning special kinds of models to count and also relations other than isomorphisms between models. Vaught's conjecture can be translated to counting the number of orbits corresponding to the action of S_{∞} , the symmetric group of ω , on the Polish space of countable models. One way to obtain a positive result is to consider only isomorphisms induced by a *subgroup G* of S_{∞} Vaught's conjecture has been confirmed when *G* is solvable; the best result in this type of investigations, is the case when *G* is a lie group.

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Our work here is inspired by Gabor Sági, who approached Vaught's conjecture using the machinery of algebraic logic.

Cylindric algebras-reflecting both syntax and semantics

For a set $V, \mathcal{B}(V)$ denotes the Boolean set algebra $\langle \wp(V), \cup, \cap, \sim, \emptyset, V \rangle$. Let U be a set and α an ordinal; α will be the dimension of the algebra. For $X \subseteq {}^{\alpha}U$ and $i, j < \alpha$, let

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The algebra $\langle \mathcal{B}(^{\alpha}U), C_i, D_{ij} \rangle_{i,j < \alpha}$ is called *the full cylindric set algebra of dimension* α with unit (or greatest element) $^{\alpha}U$. Any subalgebra of the latter is called a set algebra of dimension α .

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Then the set $\{\phi^{\mathsf{M}} : \phi \in Fm^{L}\}$ is a cylindric set algebra of dimension α , where Fm^{L} denotes the set of first order formulas taken in the signature *L*. To see why, we have:

$$\begin{split} \phi^{\mathsf{M}} \cap \psi^{\mathsf{M}} &= (\phi \land \psi)^{\mathsf{M}}, \\ {}^{\alpha}\mathsf{M} \sim \phi^{\mathsf{M}} &= (\neg \phi)^{\mathsf{M}}, \\ \mathsf{C}_{i}(\phi^{\mathsf{M}}) &= (\exists v_{i}\phi)^{\mathsf{M}}, \\ \mathsf{D}_{ij} &= (x_{i} = x_{j})^{\mathsf{M}}. \end{split}$$

By Cs_{α} we denote the class of all subalgebras of full set algebras of dimension α . The (equationally defined) **CA**_{α} class is obtained from cylindric set algebras by a process of abstraction and is defined by a *finite* schema of equations that holds of course in the more concrete set algebras.

Let α be an ordinal. By a cylindric algebra of dimension α , briefly a **CA**_{α}, we mean an algebra

$$\mathfrak{A} = \langle \mathsf{A}, +, \cdot, -, \mathsf{0}, \mathsf{1}, \mathsf{c}_i, \mathsf{d}_{ij} \rangle_{\kappa, \lambda < lpha}$$

where $\langle A, +, \cdot, -, 0, 1 \rangle$ is a Boolean algebra such that 0, 1, and d_{ij} are distinguished elements of A (for all $j, i < \alpha$), - and c_i are unary operations on A (for all $i < \alpha$), + and . are binary operations on A, and such that the following equations are satisfied for any $x, y \in A$ and any $i, j, \mu < \alpha$:

$$\begin{array}{l} (C_1) \ c_i 0 = 0, \\ (C_2) \ x \leq c_i x \ (i.e., x + c_i x = c_i x), \\ (C_3) \ c_i (x \cdot c_i y) = c_i x \cdot c_i y, \\ (C_4) \ c_i c_j x = c_j c_i x, \\ (C_5) \ d_{ii} = 1, \\ (C_6) \ \text{if } i \neq j, \mu, \text{ then } d_{j\mu} = c_i (d_{ji} \cdot d_{i\mu}), \\ (C_7) \ \text{if } i \neq j, \text{ then } c_i (d_{ij} \cdot x) \cdot c_i (d_{ij} \cdot -x) = 0. \end{array}$$

Cylindric algebras-reflecting both syntax and semantics

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Let α be an ordinal. An algebra $\mathfrak{A} \in CA_{\alpha}$ is *locally finite*, if the dimension set of every element $x \in A$ is finite. The dimension set of x, or Δx for short, is the set $\{i \in \alpha : c_i x \neq x\}$. Locally finite algebras correspond to Tarski–Lindenbaum algebras of (first order) formulas; in such algebras the dimension set of (an equivalence class of) a formula reflects the number of (finite) set of free variables in this formula.

Tarski proved that every locally finite α -dimensional cylindric algebra is representable, i.e. isomorphic to a subdirect product of set algebra each of dimension α . Let Lf_{α} denote the class of locally finite cylindric algebras. Tarski proved that every locally finite α -dimensional cylindric algebra is representable, i.e. isomorphic to a subdirect product of set algebra each of dimension α . Let Lf_{α} denote the class of locally finite cylindric algebras.

Let \mathbf{RCA}_{α} stand for the class of isomorphic copies of subdirect products of set algebras each of dimension α , or briefly, the class of α dimensional representable cylindric algebras. Then Tarski's theorem reads $\mathbf{Lf}_{\alpha} \subseteq \mathbf{RCA}_{\alpha}$. This representation theorem is non-trivial; in fact it is equivalent to Gödel's celebrated Completeness Theorem. Completeness in the general case is a huge subject that has provoked extensive research. A natural generalization of \mathbf{Lf}_{α} is \mathbf{Dc}_{α} when α is infinite; $\mathfrak{A} \in \mathbf{Dc}_{\alpha}$ iff $\alpha \sim \Delta x$ is infinite for all $x \in A$.

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Definition:

1. A subset $X \subseteq \mathbb{R}$ is *meager* if it is a countable union of nowhere dense sets. Let covK be the least cardinal κ such that \mathbb{R} can be covered by κ many nowhere dense sets. Let \mathfrak{p} be the least cardinal κ such that there are κ many meager sets of \mathbb{R} whose union is not meager.

2. A *Polish space* is a topological space that is metrizable by a complete separable metric.

Examples of Polish spaces are \mathbb{R} , the Cantor set $^{\omega}2$ and the Baire space $^{\omega}\omega$. These are called *real spaces* because they are Baire isomorphic. Any second countable compact Hausdorff space, like the Stone space of a countable Boolean algebra, is a Polish space (a complete separable metric space).

Theorem

- 1. The cardinals covK and $\mathfrak p$ are uncountable cardinals, such that $\mathfrak p \leq \text{covK} \leq 2^\omega.$
- 2. The cardinal covK is the least cardinal such the Baire category theorem for Polish spaces fails, and it is also the largest for which Martin's axiom for countable Boolean algebras holds.
- 3. If X is a Polish space, then it cannot be covered by < covK many meager sets. If $\lambda < \mathfrak{p}$, and $(A_i : i < \lambda)$ is a family of meager subsets of X, then $\bigcup_{i \in \lambda} A_i$ is meager.

Both cardinals covK and \mathfrak{p} have an extensive literature. It is consistent that $\omega < \mathfrak{p} < \operatorname{covK} \leq 2^{\omega}$ so that the two cardinals are generally different, but it is also consistent that they are equal; equality holds for example in the Cohen real model of Solovay and Cohen. In this case, Martin's axiom implies that they are both equal to the continuum. Let \mathfrak{A} be any Boolean algebra. The set of ultrafilters of \mathfrak{A} is denoted by $\mathfrak{U}(\mathfrak{A})$. The Stone topology makes $\mathfrak{U}(\mathfrak{A})$ a compact Hausdorff space. We denote this space by \mathfrak{A}^* . Recall that the Stone topology has as its basic open sets the sets $\{N_x : x \in A\}$ where

$$N_{x} = \{F \in \mathfrak{U}(\mathfrak{A}) : x \in F\}.$$

Let $x \in A$, $Y \subseteq A$ and suppose that $x = \sum Y$. We say that an ultrafilter $F \in \mathfrak{U}(\mathfrak{A})$ preserves $Y \iff$ whenever $x \in F$, then $y \in F$ for some $y \in Y$. Now let $\mathfrak{A} \in Lf_{\omega}$. For each $i \in \omega$ and $x \in A$ let

$$\mathfrak{U}_{i,x} = \{F \in \mathfrak{U}(\mathfrak{A}) : F \text{ preserves } \{s_i^j x : j \in \omega\}\}.$$

Then

$$\begin{split} \mathfrak{U}_{i,x} &= \{F \in \mathfrak{U}(\mathfrak{A}) : \mathsf{c}_i x \in F \Rightarrow (\exists j \in \omega) \mathsf{s}_j^i x \in F\} \\ &= \mathsf{N}_{-\mathsf{c}_i x} \cup \bigcup_{j < \omega} \mathsf{N}_{\mathsf{s}_j^j x}. \end{split}$$

Counting models omitting types for quantifier logics with infinitely many variables

Let

$$\mathcal{H}(\mathfrak{A}) = \bigcap_{i \in \omega, x \in A} \mathfrak{U}_{i,x}(\mathfrak{A}) \cap \bigcap_{i \neq j} N_{-\mathsf{d}_{ij}}.$$

It is clear that $\mathcal{H}(\mathfrak{A})$ is a G_{δ} set in \mathfrak{A}^* . For $F \in \mathfrak{U}(\mathfrak{A})$, let

$$rep_F(x) = \{ \tau \in {}^{\omega}\omega : \mathsf{s}^{\mathfrak{A}}_{\tau}x \in F \},\$$

for all $x \in A$. Here for $\tau \in {}^{\omega}\omega$, $s_{\tau}^{\mathfrak{A}}x$ by definition is $s_{\tau \upharpoonright \Delta x}^{\mathfrak{A}}x$. The latter is well defined because $|\Delta x| < \omega$. When $a \in F$, then rep_F is a representation of \mathfrak{A} such that $rep_F(a) \neq 0$.

The following theorem establishes a one to one correspondence between representations of locally finite cylindric algebras and Henkin ultrafilters. Cs_{ω}^{reg} denotes the class of *regular set algebras*; a a set algebra with top element ${}^{\alpha}U$ is such, if whenever $f, g \in {}^{\alpha}U, f \upharpoonright \Delta x = g \upharpoonright \Delta x$, and $f \in X$ then $g \in X$. This reflects the metalogical property that if two assignments agree on the free variables occurring in a formula then both satisfy the formula or none does.

If $F \in \mathcal{H}(\mathfrak{A})$, then rep_F is a homomorphism from \mathfrak{A} onto an element of $Lf_{\omega} \cap Cs_{\omega}^{reg}$ with base ω . Conversely, if h is a homomorphism from \mathfrak{A} onto an element of $Lf_{\omega} \cap Cs_{\omega}^{reg}$ with base ω , then there is a unique $F \in \mathcal{H}(\mathfrak{A})$ such that $h = rep_F$.

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Lemma

Suppose that *T* is a theory, $|T| = \lambda$, λ regular, then there exist models \mathfrak{M}_i : $i < \lambda^2$, each of cardinality λ , such that if $i(1) \neq i(2) < \chi$, $\bar{a}_{i(l)} \in M_{i(l)}$, l = 1, 2,, $\operatorname{tp}(\bar{a}_{l(1)}) = \operatorname{tp}(\bar{a}_{l(2)})$, then there are $p_i \subseteq \operatorname{tp}(\bar{a}_{l(i)})$, $|p_i| < \lambda$ and $p_i \vdash \operatorname{tp}(\bar{a}_{l(i)})$ (tp (\bar{a}) denotes the complete type realized by the tuple \bar{a}). We shall use the algebraic counterpart of the following corollary obtained by restricting Shelah's theorem to the countable case:

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Corollary

For any countable theory, there is a family of $< {}^{\omega}2$ countable models that overlap only on principal types.

Theorem

Assume that $\kappa < \mathfrak{p}$. Let α be a countable infinite ordinal.

- 1. Let $\mathfrak{A} \in Dc_{\alpha}$ be countable. Let $(\Gamma_i : i \in \kappa)$ be a set of non-principal types in \mathfrak{A} . Then there is a weak set algebra \mathfrak{B} , that is, \mathfrak{B} has top element a weak space, and a homomorphism $f : \mathfrak{A} \to \mathfrak{B}$ such that for all $i \in \kappa$, $\bigcap_{x \in X_i} f(x) = \emptyset$, and $f(a) \neq 0$. If \mathfrak{A} is simple, then \mathfrak{p} can be replaced by covK.
- 2. If $\mathfrak{A} \in Lf_{\alpha}$, and $(\Gamma_i : i \in \kappa)$ is a family of finitary non-principal types then there is a topological set algebra \mathfrak{B} , that is, \mathfrak{B} has top element a Cartesian square, and $\mathfrak{B} \in Cs_{\alpha}^{reg} \cap Lf_{\alpha}$ together with a homomorphism $f : \mathfrak{A} \to \mathfrak{B}$ such that $\bigcap_{x \in X_i} f(x) = \emptyset$, and $f(a) \neq 0$. If the family of given types are ultrafilters then \mathfrak{p} can be replaced by 2^{ω} , so that $< 2^{\omega}$ types can be omitted.

Proof

For the first part, we have

$$(\forall j < \alpha)(\forall x \in A)(c_j x = \sum_{i \in \alpha \smallsetminus \Delta x} s_i^j x.)$$
 (1)

Now let *V* be the weak space ${}^{\omega}\omega^{(ld)} = \{s \in {}^{\omega}\omega : |\{i \in \omega : s_i \neq i\}| < \omega\}$. For each $\tau \in V$ for each $i \in \kappa$, let

$$X_{i,\tau} = \{\mathbf{s}_{\tau}\mathbf{x} : \mathbf{x} \in X_i\}.$$

Here s_{τ} is the unary operation as defined corresponding to τ . For each $\tau \in V$, s_{τ} is a complete Boolean endomorphism on \mathfrak{A} by It thus follows that

$$(\forall \tau \in V)(\forall i \in \kappa) \prod^{\mathfrak{A}} X_{i,\tau} = 0$$
 (2)

Counting models omitting types for quantifier logics with infinitely many variables

Proof (Continue)

Let *S* be the Stone space of the Boolean part of \mathfrak{A} , and for $x \in \mathfrak{A}$, let N_x denote the clopen set consisting of all Boolean ultrafilters that contain *x*. Then from (1) and (2) it follows that for $x \in \mathfrak{A}$, $j < \beta$, $i < \kappa$ and $\tau \in V$, the sets

$$\mathbf{G}_{j,x} = N_{\mathsf{c}_{j}x} \setminus igcup_{i
otin \Delta x} N_{\mathsf{s}_{j}^{i}x} ext{ and } \mathbf{H}_{i, au} = igcap_{x \in X_{i}} N_{\mathsf{s}_{ar{ au}}x}$$

are closed nowhere dense sets in S. Also each $\mathbf{H}_{i,\tau}$ is closed and nowhere dense. Let

$$\mathbf{G} = \bigcup_{j \in \beta} \bigcup_{x \in B} \mathbf{G}_{j,x} \text{ and } \mathbf{H} = \bigcup_{i \in \kappa} \bigcup_{\tau \in V} \mathbf{H}_{i,\tau}.$$

By properties of \mathfrak{p} , **H** can be reduced to a countable collection of nowhere dense sets.

Proof (Continue)

By the Baire Category theorem for compact Hausdorff spaces, we get that $\mathfrak{H}(\mathfrak{A}) = S \sim \mathbf{H} \cup \mathbf{G}$ is dense in S. Let F be an ultrafilter in $N_a \cap X$. By the very choice of F, it follows that $a \in F$ and we have the following.

$$(\forall j < \beta)(\forall x \in B)(\mathbf{c}_j x \in F) \implies (\exists j \notin \Delta x) \mathbf{s}_j^i x \in F.)$$

$$(3)$$

and

$$(\forall i < \kappa)(\forall \tau \in V)(\exists x \in X_i)s_{\tau}x \notin F.$$
(4)

Let $V = {}^{\omega}\omega^{(d)}$ and let W be the quotient of V as defined above. That is $W = V/\bar{E}$ where $\tau \bar{E} \sigma$ if $d_{\tau(i),\sigma(i)} \in F$ for all $i \in \omega$.

Counting models omitting types for quantifier logics with infinitely many variables

Define f by $f(x) = \{\bar{\tau} \in W : s_{\tau}x \in F\}$, for $x \in \mathfrak{A}$. Then f is a homomorphism such that $f(a) \neq 0$ and it can be easily checked that $\bigcap f(X_i) = \emptyset$ for all $i \in \kappa$, hence the desired conclusion. If \mathfrak{A} is simple, then by the properties of covK, $\mathfrak{H}(\mathfrak{A}) = S \sim \mathbf{H} \cup \mathbf{G}$ is non-empty. Let $F \in H(\mathfrak{A})$ and let $a \in F$. The representation built using such F as above, call it f, has $f(a) \neq 0$, By simplicity of \mathfrak{A} , f is an injection, because $kerf = \{0\}$, since $a \notin kerf$ and by simplicity, either $kerf = \{0\}$ or $kerf = \mathfrak{A}$. 2. One proceeds exactly like in the previous item, but using, as indicated above, the fact that the operations s_{τ} for $any \tau \in {}^{\omega}\omega$ which are definable in locally finite algebras, via $s_{\tau}x = s_{\tau \upharpoonright \Delta x}x$, for any $x \in A$. Furthermore, $s_{\tau} \upharpoonright \mathfrak{Nr}_n \mathfrak{A}$ is a complete Boolean endomorphism, so that we guarantee that infimums are preserved and the sets $\mathbf{H}_{i,\tau} = \bigcap_{x \in X_i} N_{s_{\tau}x}$ remain no-where dense in the Stone topology.

2. One proceeds exactly like in the previous item, but using, as indicated above, the fact that the operations s_{τ} for $any \tau \in {}^{\omega}\omega$ which are definable in locally finite algebras, via $s_{\tau}x = s_{\tau \upharpoonright \Delta x}x$, for any $x \in A$. Furthermore, $s_{\tau} \upharpoonright \mathfrak{Mr}_n \mathfrak{A}$ is a complete Boolean endomorphism, so that we guarantee that infimums are preserved and the sets $\mathbf{H}_{i,\tau} = \bigcap_{x \in X_i} N_{s_{\tau}x}$ remain no-where dense in the Stone topology.

Now for the second part. Let $\mathfrak{A} \in Lf_{\alpha}$, $\lambda < 2^{\omega}$ and $\mathbf{F} = (X_i : i < \lambda)$ be a family of maximal non-principal finitary types, so that for each $i < \lambda$, there exists $n \in \omega$ such that $X_i \subseteq \mathfrak{Mr}_n\mathfrak{A}$, and $\prod X_i = 0$; that is X_i is a Boolean ultrafilter in $\mathfrak{Mr}_n\mathfrak{A}$.

Then by Theorem 1, or rather its direct algebraic counterpart, there are ${}^{\omega}2$ representations such that if X is an ultrafilter in $\mathfrak{Mr}_n\mathfrak{A}$ (some $n \in \omega$)) that is realized in two such representations, then X is necessarily principal. That is there exist a family of countable locally finite set algebras, each with countable base, call it $(\mathfrak{B}_{j_i} : i < 2^{\omega})$, and isomorphisms $f_i : \mathfrak{A} \to \mathfrak{B}_{j_i}$ such that if X is an ultrafilter in $\mathfrak{Mr}_n\mathfrak{A}$, for which there exists distinct $k, l \in 2^{\omega}$ with $\bigcap f_i(X) \neq \emptyset$ and $\bigcap f_j(X) \neq \emptyset$, then X is principal, so that from Shelah's lemma such representations overlap only on maximal principal types.

Then there exists a family $(F_i : i < 2^{\omega})$ of Henkin ultrafilters such that $f_i = h_{F_i}$, and we can assume that h_{F_i} is an \mathbf{CA}_{α} isomorphism as follows. Denote F_i by G. Assume, for contradiction, that there is no representation (model) that omits \mathbf{F} . Then for all $i < 2^{\omega}$, there exists F such that F is realized in \mathfrak{B}_{j_i} . Let $\psi : 2^{\omega} \to \wp(\mathbf{F})$, be defined by $\psi(i) = \{F : F \text{ is realized in } \mathfrak{B}_{j_i}\}$. Then for all $i < 2^{\omega}$, $\psi(i) \neq \emptyset$. Furthermore, for $i \neq k$, $\psi(i) \cap \psi(k) = \emptyset$, for if $F \in \psi(i) \cap \psi(k)$ then it will be realized in \mathfrak{B}_{j_i} and \mathfrak{B}_{j_k} , and so it will be principal. This implies that $|\mathbf{F}| = 2^{\omega}$ which is impossible.

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The Silver Vaught Dichotomy asserts that there are either countably many equivalence classes or there is a perfect set of pairwise inequivalent elements. For any continous action by a Polish group G on a Polish space X, the orbit equivalence relation is conjectured to satisfy the Silver Vaught Dichotomy. This conjecture both implies and is motivated by Vaught's conjecture. In Vaught's conjecture is the particular case. when the group is the symmetric group of permutations on X, and the set X, is the set of non isomorphic models of a theory with domain ω . The relation E is just the equivalence relation of isomorphism. In our case X was a G_{δ} subset of the Stone space of a countable cylindric algebra.

Given an equivalence relation there are theorems that assert that either the quotient space is 'small' or else it contains a copy of a specific 'large' set. Two dichotomies showing this tendency are known.

Another Dichotomy, called the Glimm Effros dchotomy for an equivalence relation E asserts that E contains a copy of the Vitali equivalence relation E_0 (equivalently there exists a non atomic ergodic measure for E) or else there is a countable Borel separating family for E. This dichotomy originates with Glimm and Effros and is motivated by guestions about operator algebra. Glimm proved that the orbit space of a Polish group G action satisfies the Glimm Effros Dichotomy if G is locally compact. Effros proved it for any Polish group G, provded that the equivalence relation is F_{σ} . There exists S_{∞} spaces which violate the Glimm Effros dichotomy, but for the Silver Vaught Dichotomy this is still an open question. In all cases we consider, the Glimm Effros dichotomy implies the Silver Vaught dichotomy.

Counting models omitting types for quantifier logics with infinitely many variables

Theorem

Let $G \subseteq S_{\infty}$ be a cli group, and let E_G denote the corresponding orbit equivalence relation. Then $|\mathcal{H}(\mathfrak{A})/E_G| \leq \omega$ or $|\mathcal{H}(\mathfrak{A})/E_G| = 2^{\omega}$

Let $G \subseteq S_{\infty}$ be a cli group, and let E_G denote the corresponding orbit equivalence relation. Then $|\mathcal{H}(\mathfrak{A})/E_G| \leq \omega$ or $|\mathcal{H}(\mathfrak{A})/E_G| = 2^{\omega}$

Proof

It is known that the number of orbits of E_G satisfies the so-called Glimm-Effros Dichotomy. By known results in the literature on the topological version of Vaught's conjecture, we have $\mathcal{H}(\mathfrak{A})/E_G$ is either at most countable or $\mathcal{H}(\mathfrak{A})/E_G$ contains continuum many non equivalent elements (i.e non-isomorphic models).

It is known that the number of orbits of $E = E_{S_{\infty}}$ does not satisfy the Glimm Effros Dichotomy. We note that cli groups cover all natural extensions of abelian groups, like nilpotent and solvable groups. Now we give a topological condition that implies Vaught's conjecture. Let everything be as above with *G* denoting a Polish subgroup of S_{∞} . Give $\mathcal{H}(\mathfrak{A})/E_G$ the quotient topology and let $\pi : \mathcal{H}(\mathfrak{A}) \to \mathcal{H}(\mathfrak{A})/E_G$ be the projection map. π of course depends on *G*, we sometimes denote it by π_G to emphasize the dependence.

Lemma

 π is open.

Lemma

 π is open.

Proof

To show that π is open it is enough to show for arbitrary $a \in \mathfrak{A}$ that $\pi^{-1}(\pi(N_a))$ is open. For,

$$\pi^{-1}(\pi(N_{a})) = \{F \in \mathcal{H}(\mathfrak{A}) : (\exists F' \in N_{a}) (F, F') \in E\}$$

$$= \{F \in \mathcal{H}(\mathfrak{A}) : (\exists F' \in N_{a})(\exists \rho \in G) \ s_{\rho}^{+}F' = F\}$$

$$= \{F \in \mathcal{H}(\mathfrak{A}) : (\exists F' \in N_{a})(\exists \rho \in G) \ F' = s_{\rho^{-1}}^{+}F\}$$

$$= \{F \in \mathcal{H}(\mathfrak{A}) : (\exists \rho \in G)s_{\rho^{-1}}^{+}F \in N_{a}\}$$

$$= \{F \in \mathcal{H}(\mathfrak{A}) : (\exists \rho \in G)a \in s_{\rho^{-1}}^{+}F\}$$

$$= \{F \in \mathcal{H}(\mathfrak{A}) : (\exists \rho \in G)s_{\rho}^{+}a \in F\}$$

$$= \bigcup_{\rho \in G} N_{s_{\rho}^{+}a}$$

Counting models omitting types for quantifier logics with infinitely many variables

Theorem

If π is closed, then Vaught's conjecture holds.

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Proof

We have $\mathcal{H}(\mathfrak{A})$ is Borel subset of \mathfrak{A}^* , the Stone space of \mathfrak{A} , and $\mathcal{H}(\mathfrak{A})/E_G$ is a continuous image of $\mathcal{H}(\mathfrak{A})$. Because π is open, $\mathcal{H}(\mathfrak{A})/E_G$ is second countable. Now, since $\mathcal{H}(\mathfrak{A})$ is metrizable, it is normal. Since π is closed, open, continuous, and surjective, so $\mathcal{H}(\mathfrak{A})/E_G$ is also normal, hence regular. Thus $\mathcal{H}(\mathfrak{A})/E_G$ can be embedded in \mathbb{R}^{ω} (like in the proof of Urysohn's metrization Theorem). If $\mathcal{H}(\mathfrak{A})/E_G$ is uncountable, then being analytic (the continuous image under a map between two Polish spaces of a Borel set), it has the power of the continuum. Unfortunately, π can't be closed when $G = S_{\infty}$ (or *G* sufficiently large as we shall see) and \mathfrak{A} is simple (this is the case when our theory *T* is complete). Indeed, if it was closed, then as has just been shown, $\mathcal{H}(\mathfrak{A})/E$ is Haussdorf. A well known fact says that: when the quotient map is open, $\mathcal{H}(\mathfrak{A})/E$ is Haussdorf iff *E* is closed. We show that when \mathfrak{A} is simple, then *E* is not closed. For, assume towards a contradiction that *E* is closed, that is $\sim E$ is open. Let $(F, F') \notin E$. Then for some $a \in F, b \in F', N_a \times N_b \cap E = \emptyset$, i.e., for all $\tau \in S_{\infty}, a.s_{\tau}^+b = 0$. This last situation is of course impossible because one can choose τ so that $\Delta a \cap \Delta s_{\tau}^+b = \emptyset$. Here we we used the fact that when \mathfrak{A} is simple and $\Delta x \cap \Delta y = \emptyset$, then $x.y \neq 0$.

An algebraic proof to Morley's theorem endowed with OTT

We next give a new proof of Morley's theorem; we also count the number of models omitting a given family of types.

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Theorem

Suppose T is a first order complete theory in a countable language with equality.

- 1. (Morley) If T has more than ω_1 countable models, then it has 2^{ω} countable models. The same statement holds for theories not necessarily complete, in countable languages with or without equality.
- 2. If $(\Gamma_i : i < \omega)$ be a family of non-isolated types, then the number of non isomorphic countable models, omitting this family, is either ω , ω_1 or $^{\omega}$ 2

Proof

Let *T* be a first order theory in a countable language with equality, and let $\mathfrak{A} = \mathfrak{Fm}_{T}$. Then S_{∞} is a Polish group with respect to composition of functions and the topology it inherits from the Baire space ${}^{\omega}\omega$. Consider the map $J : S_{\infty} \times \mathcal{H}(\mathfrak{A}) \longrightarrow \mathcal{H}(\mathfrak{A})$ defined by $J(\rho, F) = s_{\rho}^{+}F$ for all $\rho \in S_{\infty}$, $F \in \mathcal{H}(\mathfrak{A})$. Then *J* is a well defined action of S_{∞} on $\mathcal{H}(\mathfrak{A})$. Also *J* is a continuous map from $S_{\infty} \times \mathcal{H}(\mathfrak{A})$ (with the product topology) to $\mathcal{H}(\mathfrak{A})$ because for an arbitrary $a \in A$,

$$J^{-1}(N_a \cap \mathcal{H}(\mathfrak{A})) = \bigcup_{\tau \in \mathsf{S}_{\infty}} (\{\mu^{-1} : \mu \in \mathsf{S}_{\infty}, \, \mu|_{\Delta a})$$

 $=\tau|_{\Delta a}\}\times[N_{\mathbf{s}_{\tau}^{+}a}\cap\mathcal{H}(\mathfrak{A})]).$

It follows that the the orbit equivalence relation is analytic. By Burgess' Theorem if there are more than ω_1 orbits, then there are 2^{ω} orbits. But the number of orbits here is exactly the number of non-isomorphic countably infinite models of \mathcal{T} . This completes the proof. For the part on omitting types, set $\mathcal{H}_{omit} = \mathcal{H}(\mathfrak{Fm}_{\mathcal{T}}) \cap \bigcap_{i \in \omega, \tau \in W} \bigcup_{\varphi \in \Gamma_i} N_{-s_{\tau}^+(\varphi/\equiv_{\tau})}$, where $W = \{\tau \in {}^{\omega}\omega : |i : \tau(i) \neq i| < \omega\}$. Clearly, $|\mathcal{H}_{omit}$ is G_{δ} , so it is Polish. For the remaining part one uses locally finite QA_ws instead of Lf_ws.

Let *T* be a countable theory. Then the number of non isomorphic models is equal to the number of models omitting a given a set of $< \lambda$ many types are the same $\iff |\mathcal{H}(\mathfrak{Fm}_{T})| > |\bigcup_{i \in \lambda, \tau \in W} \bigcap_{\varphi \in \Gamma_{i}} N_{s_{\tau}^{+}(\varphi/\equiv \tau)}|.$ Let T be a countable theory. Then the number of non isomorphic models is equal to the number of models omitting a given a set of $< \lambda$ many types are the same $\iff |\mathcal{H}(\mathfrak{Fm}_{T})| > |\bigcup_{i \in \lambda, \tau \in W} \bigcap_{\varphi \in \Gamma_{i}} N_{s_{\tau}^{+}(\varphi/\equiv \tau)}|.$

The next example shows that this may fail to happen: Consider non standard models of arithmetic. \mathbb{N} is an atomic model, which means that the neat *n*-reduct of the locally finite cylindric algebra \mathfrak{Fm}_T based on $T = \text{Th}(\mathbb{N})$ is atomic for each *n*. For each $n \in \omega$, let Γ_n be the set of co-atoms in the neat *n*-reduct. These are non-principal types and a model M omits them \iff it is atomic, hence it is isomorphic to \mathbb{N} because atomic models are unique. But Peano arithmetic is unstable, so it has ω^2 many non-isomorphic countable models (non-standard models of arithmetic).

Another example exhibiting the same phenomena: Let T be the theory of algebraically closed fields of characteristic zero. Then T is ω stable and it has countably many non-isomorphic models; for each $\alpha \leq \omega$, there is a model of transcendence degree α over the rationals. Take the types as above. In this case the all subalgebras of the *n*-neat reducts are atomic. Then the the field of algebraic number is the only countable model omitting this family of types. This is an atomic model. This theory has also another countable ω -saturated model, which is that of transcendence degree ω .

Example 2

There is a somewhat amusing Theorem of Vaught's that says that a countable theory cannot have exactly two models. We show that this is not the case when we require that the constructed odels omit a given family of non-principal types. Take the language $L = \{c_n : n \in \omega\}$. Then a model M of T is determined by how many extra elements it has, i.e by $|\{b \in M : b \neq c_n^M\}|$. So T is ω_1 categorical and since T has only infinite models it is complete. Also T has countably many non isomorphic models, M_{α} with α extra elements for $\alpha \leq \omega$. Consider the *m* type $\Gamma = \bigwedge_{i \neq i < m} \{ v_i \neq v_j \} \cup \{ v_0 \neq c_n : n \in I \}$ ω \cup { $v_1 \neq c_n : n \in \omega$ } ... { $v_{m-1} \neq c_n : n \in \omega$ }. Then Γ is non-principal and it is omitted by exactly m models namely $M_0, M_1, \ldots, M_{m-1}$. This can be generalized for complete strongly minimal theories which have countable models of dimension α , $\alpha < \omega$.

Example 3

We show that there is a theory having exactly ω_1 models omitting continuum many types. Take the first order countable theory in the language $\{<, c_0, c_1, \dots\}$ where < is a binary relation symbols and the c'_i 's ($i \le \omega$) are constants. Let T be the L theory which states that < is a linear order and that $c_i \neq c_i$ for $i \neq j$. Take $\Gamma_1 = \{v_1 \neq c_i : i \in \omega\}$ and for every injective $f \in {}^{\omega}\omega$, let $\Gamma_f = \{c_{f(i)} > c_{f(i+1)} : i \in \omega\}$. Consider the set of non-principal types $\mathcal{G} = \{\Gamma_1, \Gamma_f : f \in \widehat{\omega}\omega\}$. Then a model M omits $\mathcal{G} \iff$ it is a countable well order. The family \mathcal{G} is uncountable. Making this family countable would violate Vaught's conjecture in $L_{\omega_1,\omega}$. Indeed let T be a countable theory and $\{\Gamma_i : i < \omega\}$ be a family of non-principal types omitted by exactly ω_1 models. Then the $L_{\omega_1,\omega}$ sentence $\bigwedge T \land \bigwedge_{n \in \omega} (\neg (\exists v_n) \bigwedge_{\phi \in \Gamma_n} \phi(v_n))$ violates Vaught's conjecture; for it has ω_1 countable models.

Vaught's conjecture holding for distinguishable ordinary models and pairwise non-isomorphic models

Vaught's conjecture holding for distinguishable ordinary models and pairwise non-isomorphic models

We define an equivalence relation on ultrafilters that turns out to be Borel. This implies that it satisfies the Glimm-Effros dichotomy, and so has either countably many or else continuum many equivalence classes. The equivalence relation we introduce corresponds to a non-trivial equivalence relation between models which is weaker than isomorphism and stronger than elementary equivalence. We define an equivalence relation on ultrafilters that turns out to be Borel. This implies that it satisfies the Glimm-Effros dichotomy, and so has either countably many or else continuum many equivalence classes. The equivalence relation we introduce corresponds to a non-trivial equivalence relation between models which is weaker than isomorphism and stronger than elementary equivalence.

Definition (Notation)

Let \mathcal{F} be an ultrafilter of a locally finite (cylindric or quasi-polyadic) algebra \mathfrak{A} . For $a \in A$ define

$$Sat_{\mathcal{F}}(a) = \{t|_{\Delta a} : t \in {}^{\omega}\omega, \ s_t^+a \in \mathcal{F}\}.$$

Throughout, \mathfrak{A} is countable. We define an equivalence relation \mathcal{E} on the space $\mathcal{H}(\mathfrak{A})$) that turns out to be Borel.

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Definition Let \mathcal{E} be the following equivalence relation on $\mathcal{H}(\mathfrak{A})$:

 $\mathcal{E} = \{ (\mathcal{F}_0, \mathcal{F}_1) : (\forall a \in A) (|Sat_{\mathcal{F}_0}(a)| = |Sat_{\mathcal{F}_1}(a)|) \}.$

We say that $\mathcal{F}_0, \mathcal{F}_1 \in \mathcal{H}(\mathfrak{A})$ are distinguishable if $(\mathcal{F}_0, \mathcal{F}_1) \notin \mathcal{E}$. We also say that two models of a theory T are distinguishable if their corresponding ultrafilters in $\mathcal{H}(CA(T) = \operatorname{Fm}_T)$ are distinguishable. That is, two models are distinguishable if they disagree in the number of realizations they have for some formula. Then \mathcal{E} is Borel in the product space $\mathcal{H}(\mathfrak{A}) \times \mathcal{H}(\mathfrak{A})$. If X be a Polish space and E a Borel equivalence relation on X. We call Esmooth if there is a Borel map f from X to the Cantor space \mathfrak{L}^2 such that

$$xEy \Leftrightarrow f(x) = f(y).$$

Note that *E* is smooth iff *E* admits a *countable Borel separating family*, i.e., a family (A_n) of Borel sets such that

$$xEy \Leftrightarrow \forall n(x \in A_n \leftrightarrow y \in A_n).$$

Clearly, if *E* is smooth then it is Borel (but the converse is not true). A standard example of a non-smooth Borel equivalence relation is the following: On $2^{\mathbb{N}}$, let E_0 be defined by

$$xE_0y \Leftrightarrow \exists n \forall m \geq n(x(m) = y(m)).$$

We say that the equivalence relation E, on a Polish space X, satisfies the *Glimm-Effros Dichotomy* if either it is smooth or else it contains a copy of E_0 . Clearly, for an equivalence relation E, E satisfies the Glimm-Effros Dichotomy implies that E satisfies the *Silver-Vaught Dichotomy*, that is, E has either countably many classes or else perfectly many classes (X has a perfect subset of non-equivalent elements).

Theorem (Harrington-Kechris-Louveau)

Let X be a Polish space and E a Borel equivalence relation on X. Then E satisfies the Glimm-Effros Dichotomy.

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It follows directly from this theorem, replacing X with $\mathcal{H}(\mathfrak{A})$ that \mathcal{E} satisfies the Glimm-Effros dichotomy and so has either countably many equivalence classes or else perfectly many.

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Corollary

Let T be a first order theory in a countable language (with or without equality). If T has an uncountable set of countable models that are pairwise distinguishable, then actually it has such a set of size 2^{\aleph_0} .

Vaught's conjecture holds when counting weak models in rich languages

Vaught's conjecture holds when counting weak models in rich languages

A rich language is one for which outside any atomic formula there are infinitely many variables–and the rest is like first order logic. Recall that rich languages (corresponding to Dc_{α}) enjoy an omitting types theorem; for < p many non-principal types, and the types can contain infinitely many variables (unlike first order logic). However, the models that omit a countable set of non-principal types is only a weak model, and it can be proved that there are cases, where it has to be a weak model.

Let *T* be the theory of dense linear order without endpoints. Then *T* is complete. Let $\Gamma(x_0, x_1 \dots)$ be the set

$$\{x_1 < x_0, x_2 < x_1, x_3 < x_2 \ldots\}.$$

(Here there is no bound on free variables.) A model \mathfrak{M} omits Γ if and only if \mathfrak{M} is a well ordering. But *T* has no well ordered models, so no model of *T* omits Γ . However *T* locally omits Γ because if $\phi(x_0, \ldots, x_{n-1})$ is consistent with *T*, then $\phi \land \neg x_{n+2} < x_{n+1}$ is consistent with *T*. Note that Γ can be omitted in a weak model.

But first a some definitions.

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Definition

Let \mathfrak{A} and \mathfrak{B} be set algebras with bases U and W respectively. Then \mathfrak{A} and \mathfrak{B} are *base isomorphic* if there exists a bijection $f : U \to W$ such that $\overline{f} : \mathfrak{A} \to \mathfrak{B}$ defined by $\overline{f}(X) = \{y \in {}^{\alpha}W : f^{-1} \circ y \in x\}$ is an isomorphism from \mathfrak{A} to \mathfrak{B}

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Definition

An algebra \mathfrak{A} is *hereditary atomic*, if each of its subalgebras is atomic.

Finite Boolean algebras are hereditary atomic of course, but there are infinite hereditary atomic Boolean algebras; any Boolean algebra generated by by its atoms is hereditary atomic, for example the finite co-finite algebra on any set. An algebra that is infinite and complete is not hereditary atomic, whether atomic or not.

Hereditary atomic algebras arise naturally as the Tarski-Lindenbaum algebras of certain countable first order theories, that abound. If *T* is a countable complete first order theory which has an an ω -saturated model, then for each $n \in \omega$, the Tarski-Lindenbuam Boolean algebra \mathfrak{Fm}_n/T is hereditary atomic. Here \mathfrak{Fm}_n is the set of formulas using only *n* variables. For example $Th(\mathbb{Q}, <)$ is such with \mathbb{Q} the ω saturated model.

A well known model-theoretic result is that T has an ω saturated model iff T has countably many n types for all n. Algebraically n-types are just ultrafilters in \mathfrak{Fm}_n/T . And indeed, what characterizes hereditary atomic algebras is that the base of their Stone space, that is the set of all ultrafilters, is at most countable.

A well known model-theoretic result is that T has an ω saturated model iff T has countably many n types for all n. Algebraically n-types are just ultrafilters in \mathfrak{Fm}_n/T . And indeed, what characterizes hereditary atomic algebras is that the base of their Stone space, that is the set of all ultrafilters, is at most countable.

Lemma

Let \mathfrak{B} be a countable Boolean algebra. If \mathfrak{B} is hereditary atomic then the number of ultrafilters is at most countable; of course they are finite if \mathfrak{B} is finite. If \mathfrak{B} is not hereditary atomic the it has 2^{ω} ultrafilters.

Our next theorem is the natural extension of Vaught's theorem to variable rich languages. However, we address only languages with finitely many relation symbols. (Our algebras are finitely generated, and being simple, this is equivalent to that it is generated by a single element.) Our next theorem is the natural extension of Vaught's theorem to variable rich languages. However, we address only languages with finitely many relation symbols. (Our algebras are finitely generated, and being simple, this is equivalent to that it is generated by a single element.)

Theorem

Let $\mathfrak{A} \in Dc_{\alpha}$ be countable simple and finitely generated. Then the number of non-base isomorphic representations of \mathfrak{A} is 2^{ω} .

Proof

Let $V = {}^{\alpha} \alpha^{(ld)}$ and let \mathfrak{A} be as in the hypothesis. Then \mathfrak{A} cannot be atomic, least hereditary atomic. By 10, it has 2^{ω} ultrafilters. For an ultrafilter F, let $h_F(a) = \{\tau \in V : s_{\tau}a \in F\}$, $a \in \mathfrak{A}$. Then $h_F \neq 0$, indeed $Id \in h_F(a)$ for any $a \in F$, hence h_F is an injection, by simplicity of \mathfrak{A} . Now $h_F : \mathfrak{A} \to \wp(V)$; all the h_F 's have the same target algebra. We claim that $h_F(\mathfrak{A})$ is base isomorphic to $h_G(\mathfrak{A})$ iff there exists a finite bijection $\sigma \in V$ such that $s_{\sigma}F = G$. We set out to confirm our claim. Let $\sigma : \alpha \to \alpha$ be a finite bijection such that $s_{\sigma}F = G$. Define $\Psi : h_F(\mathfrak{A}) \to \wp(V)$ by $\Psi(X) = \{\tau \in V : \sigma^{-1} \circ \tau \in X\}$. Then, by definition, Ψ is a base isomorphism. We show that $\Psi(h_F(a)) = h_G(a)$ for all $a \in \mathfrak{A}$. Let $a \in A$. Let $X = \{\tau \in V : s_{\tau}a \in F\}$.

Proof (continue)

Let $Z = \Psi(X)$. Then

$$Z = \{ \tau \in V : \sigma^{-1} \circ \tau \in X \}$$

= $\{ \tau \in V : s_{\sigma^{-1} \circ \tau}(a) \in F \}$
= $\{ \tau \in V : s_{\tau}a \in s_{\sigma}F \}$
= $\{ \tau \in V : s_{\tau}a \in G \}.$
= $h_G(a)$

Conversely, assume that $\bar{\sigma}$ establishes a base isomorphism between $h_F(\mathfrak{A})$ and $h_G(\mathfrak{A})$. Then $\bar{\sigma} \circ h_F = h_G$. We show that if $a \in F$, then $s_{\sigma}a \in G$. Let $a \in F$, and let $X = h_F(a)$. Then, we have

$$\sigma \circ h_F(a) = \sigma(X)$$

= {y \in V : \sigma^{-1} \circ y \in h_F(X)}
= {y \in V : \sigma_{\sigma^{-1} \circ y} a \in F}
= h_G(a)

Now we have $h_G(a) = \{y \in V : s_y a \in G\}$. But $a \in F$. Hence $\sigma^{-1} \in h_G(a)$ so $s_{\sigma^{-1}}a \in G$, and hence $a \in s_{\sigma}G$.

Define the equivalence relation \sim on the set of ultrafilters by $F \sim G$, if there exists a finite permutation σ such that $F = s_{\sigma}G$. Then any equivalence class is countable, and so we have $^{\omega}2$ many classes, which correspond to the non base isomorphic representations of \mathfrak{A} .

Theorem

Let T be a countable theory in a rich language, with only finitely many relation symbols, and $\Gamma = \{\Gamma_i : i \in \mathfrak{p}\}$ be non isolated types. Then T has 2^{ω} weak models that omit Γ . If T is complete we can replace \mathfrak{p} by covK.

1. VC can be viewed topologically as counting the number of orbits of a (Polish) group action on a Polish space.

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Or you can change the notion of isomorphism instead of counting non-isomorphic models; we count instead the number of 'distinct' models in some sense like non elementary equivalent or something in between non isomorphic and non elementary equivalent.

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Or you can change the notion of isomorphism instead of counting non-isomorphic models; we count instead the number of 'distinct' models in some sense like non elementary equivalent or something in between non isomorphic and non elementary equivalent.

We also study the case when the group G is a Polish group where the topology is induced by a complete left invariant metric.

Main Results

Counting weak modes (defined above) and distinguishable models (defined above) satisfy VC

Counting weak modes (defined above) and distinguishable models (defined above) satisfy VC

Counting orbits of a Polish group that admits a left invariant complete metric this covers Abelian (initiated by Sami from CU) and solvable and more!

Thank You!!