

Celebrating István Németi's 80th Birthday

A conceptual-based attribute to connections between theories

Mohamed Khaled





Concept algebras ...







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A concept is

definable subset of the universe M.





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 \land , \lor , \neg



Boolean Algebras



Goerge Boole (1815 - 1864)



Boolean Algebras



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Boolean set algebras



Variety of BA's

$\mathbf{2} \stackrel{\text{\tiny def}}{=} \big\langle \{0,1\}, \wedge, \vee, \neg, 0, 1 \big\rangle$

| р | q | $p \wedge q$ |
|---|---|--------------|
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |

| р | q | $p \lor q$ |
|---|---|------------|
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 1 |

| p | $\neg p$ |
|---|----------|
| 0 | 1 |
| 1 | 0 |









M; y M; x













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 M, M^2, M^3, M^4, \ldots





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$$[v_i = v_j]^{\mathfrak{M}} = \{ \overline{a} \in M^{\omega} : a_i = a_j \} \stackrel{\text{def}}{=} D_{ij}$$








Definition

The concept algebra of ${\mathfrak M}$ is:

 $\mathfrak{Cs}(\mathfrak{M}) \stackrel{\text{\tiny def}}{=} \langle \mathit{Cs}(\mathfrak{M}), \cup, \sim, \mathit{C}_i, \mathit{D}_{ij} \rangle_{i,j < \omega}.$



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Connections between theories ...





 $\mathfrak{Cs}(\mathfrak{M}) = \mathfrak{Cs}(\mathfrak{N})$



```
\mathfrak{Cs}(\mathfrak{N}) \subseteq \mathfrak{Cs}(\mathfrak{M})
```







```
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```









Homomorphisms









A translation from the language of $\mathfrak M$ to the language of $\mathfrak N$



Homomorphisms

















 $\mathfrak{Cs}(\mathfrak{M})$

 $\mathfrak{Cs}(\mathfrak{N})$

 $\mathfrak{Cs}(\mathfrak{M}) / \theta \cong \mathfrak{Cs}(\mathfrak{N})$









$$\mathfrak{Cs}(\mathfrak{N}) \hspace{.1in} \hookrightarrow \hspace{.1in} \mathfrak{Cs}(\mathfrak{M})$$



 $\mathfrak{Cs}(\mathfrak{M}) / \theta \cong$ $\mathfrak{Cs}(\mathfrak{N})$













 \implies

Mutual Definability



Mutual Definability

 \Leftarrow



 \implies

Mutual Definability



H. Andréka, J. Madarász and I. Németi (2005)



 \implies

Mutual Definability





H. Andréka, J. Madarász and I. Németi (2005)









$$\mathfrak{Cs}(\mathfrak{M}) = \langle X \rangle$$



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Changing the language of $\mathfrak M$



$$(\mathbf{x}) = \langle X \rangle$$



Changing the language of $\mathfrak M$



$$\mathfrak{Cs}(\mathfrak{M}) = \langle X$$

Changing the language of $\mathfrak M$



Adding concepts to \mathfrak{N} to get \mathfrak{M}



$$\mathcal{RS} = \langle \mathbb{R}^4, \mathsf{col}^t \rangle$$



$$\mathcal{CS}^+ = \langle \mathbb{R}^4, \mathsf{col}^\infty, \mathsf{col}^\lambda \rangle$$



Conjecture (H. Andréka)

 $\mathfrak{Cs}(\mathcal{CS}^+) = ig\langle \mathfrak{Cs}(\mathcal{RS}), b ig
angle$

for any concept b in \mathfrak{CS}^+) not in $\mathfrak{Cs}(\mathcal{RS})$.

Examples of concept algebras ...


$\mathfrak{Nr}_{n}\mathfrak{Cs}(\mathfrak{M})$



 $\mathfrak{Nr}_{n}\mathfrak{Cs}(\mathfrak{M})$

 $\mathfrak{Nr}_1\mathfrak{Cs}(\mathfrak{M})\subseteq\cdots\subseteq\mathfrak{Nr}_n\mathfrak{Cs}(\mathfrak{M})\subseteq\cdots$











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An atom:

 $\epsilon(\sim) \land \varphi$



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The neat reduct $\mathfrak{Nr}_n\mathfrak{Cs}(\mathfrak{M})$ is atomic (in fact, finite)! An atom: $v_1 \quad v_0$ $\epsilon(\sim) \land \varphi \qquad v_2 \quad v_3$



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Equivalence relation \sim



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Homogeneous structure ${\mathfrak M}$

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 $\mathbb{Q}=\langle\mathbb{Q},<\rangle$ is homogeneous



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The 3-dimensional atoms of $\mathfrak{Cs}\mathbb{Q}$ up to symmetries of indices







• Finite Langauge!





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- Local Failure of Homogenization!





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R...





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 $R_{\bullet\bullet}$

 $\mathfrak{M} \longrightarrow \mathfrak{M}_1^+$





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 $R_{\bullet\bullet}$

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$$\mathfrak{M}\longrightarrow \mathfrak{M}_1^+\longrightarrow \mathfrak{M}_2^+$$





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$$\mathfrak{M}\longrightarrow \mathfrak{M}_1^+\longrightarrow \mathfrak{M}_2^+\cdots \longrightarrow \mathfrak{M}_{n+1}^+$$





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R...

 $\mathfrak{M} \longrightarrow \mathfrak{M}_1^+ \longrightarrow \mathfrak{M}_2^+ \cdots \longrightarrow \mathfrak{M}_{n+1}^+$

 $\mathfrak{Cs}(\mathfrak{M}) = \mathfrak{Cs}(\mathfrak{M}^+_\omega)$





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- Local Failure of Homogenization!

R...

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 $\mathfrak{Cs}(\mathfrak{M})=\mathfrak{Cs}(\mathfrak{M}^+_\omega)$

 \mathfrak{M}^+_ω is homogeneous!



Thank you!